Robustness: Prior-Independence

"Prior-independent" results give us guarantees in the event that the designer doesn't know the distribution F from which the bidders' values are drawn. In this case, we assume that their values are still drawn from a prior distribution, as in the Bayesian setting, so there is some revenue-optimal mechanism OPT(F) that we wish to approximate, we just have to do so without knowing F.

The Bulow-Klemperer Result

One famous result takes the form of resource augmentation.

Theorem 1 (Bulow Klemperer [1]). For *i.i.d.* regular single-item environments, the expected revenue of the second-price auction with n + 1 agents is at least that of the optimal auction with n agents.

Let's talk about what this theorem is saying. Instead of finding the optimal auction tailored to a distribution F for n agents, you can use the Vickrey auction, which requires no prior knowledge of the distribution, so long as:

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This result **does not hold** without these assumptions. However, it is a *very strong* result, should our setting meet these assumptions.

Proof.

The Single Sample Mechanism

Can't recruit extra buyers? Instead, we can just exclude one. This is what the single sample result says.

Theorem 2 (Dhangwatnotai Roughgarden Yan [2]). Given a random sample from a bidder's distribution, posting it as a take-it-or-leave-it price gives a $\frac{1}{2}$ -approximation to the optimal revenue.

Figure 1: Geometric intuition for a posted-price from a single sample.

Proof. In quantile space!

It turns out, using a single sample from the buyers' distribution to set reserve prices and running VCG is a good approximation to the optimal mechanism. See Hartline chapter 5 for more.

Prophet Inequalities

Summary of the setting:

- Goal: Pick one item; maximize its value.
- Gambler knows distribution for each item.
- Order is adversarial.
- Inspect each item online (see v_i) and irrevocably decide whether to take or pass forever.
- Compete with OPT = $\mathbb{E}_{\mathbf{v}}[\max_i v_i]$.

The Prophet Inequality problem was posed by Samuel-Cahn '84 [7], with the original solution and analysis that we'll see by Krengel Sucheston '78 [6] and Garling. It was brought to Algorithmic Mechanism Design by Hajiaghayi Kleinberg Sandholm '07 [3], and a new analysis for this case was developed by Kleinberg Weinberg '12 [4, 5].



Figure 2: The prophet inequality problem.

Theorem 3. There is a threshold algorithm ALG such that when the gambler takes an item if and only if its value is above T, ALG $\geq \frac{1}{2}$ OPT.

Note: Can you find two different thresholds that give this same approximation?

Proof. We consider two different ways to set the threshold. Let p denote the probability that some (at least one) $v_i \ge T$ for $i \in [n]$.

Exercise: You could see this as a mechanism for a buyer to maximize social welfare. Could you design a mechanism to maximize revenue using the prophet inequality?

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